

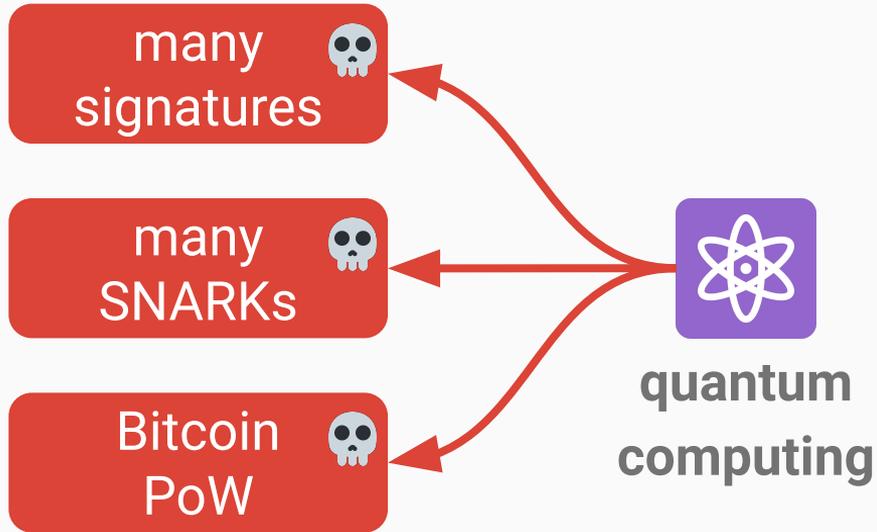


one-shot signatures

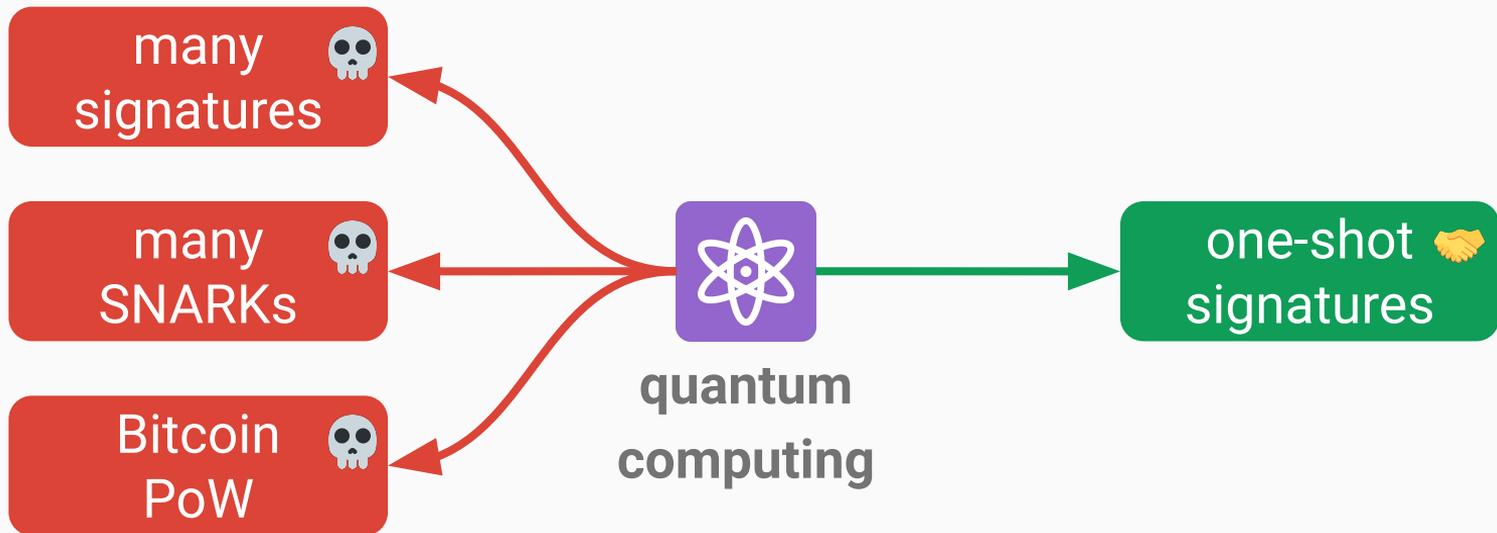
a new blockchain paradigm



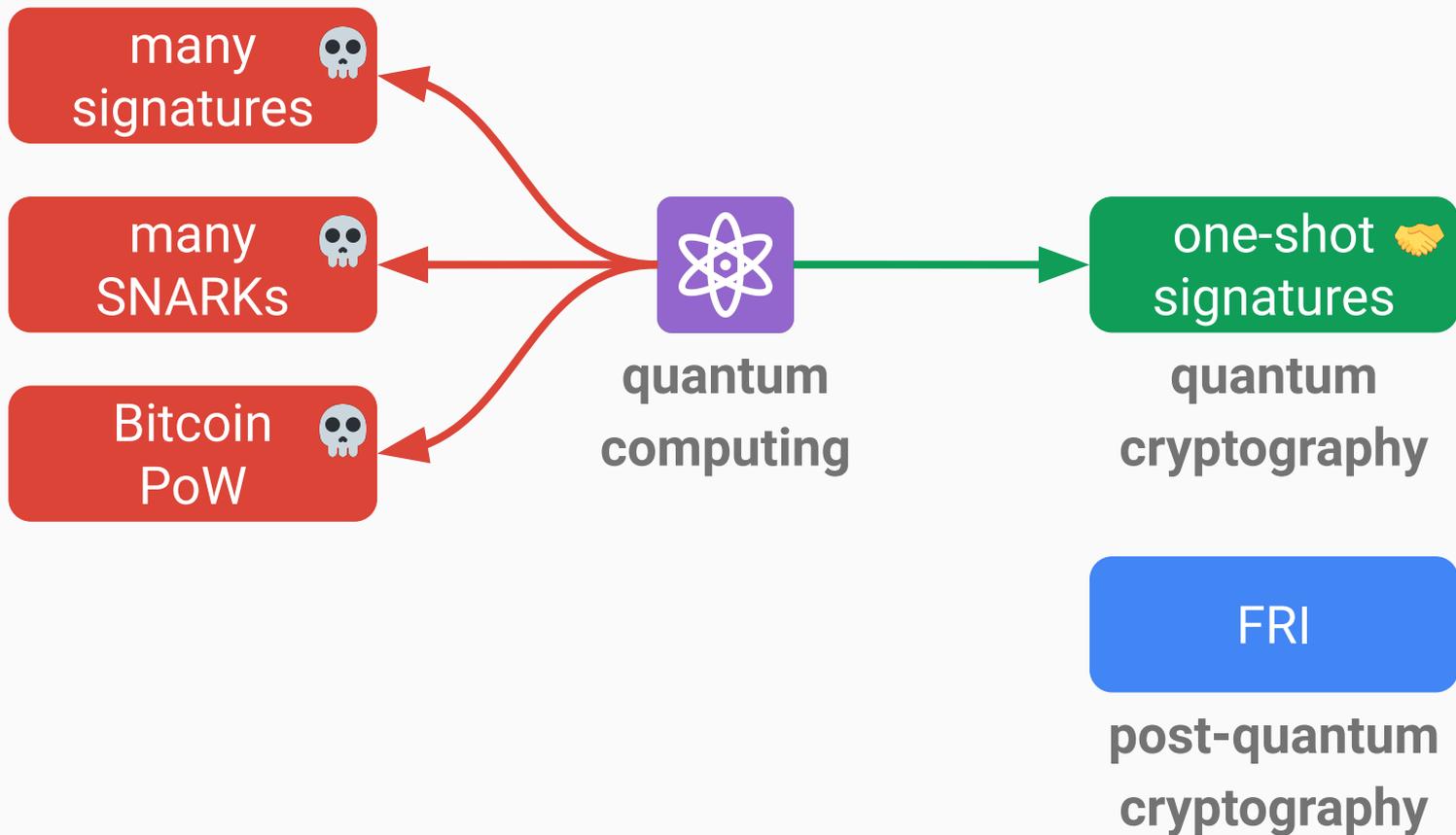
destructive vs constructive



destructive vs constructive



destructive vs constructive



part 1—signature chains

part 2—applications

part 3—accelerationism

part 1—signature chains

part 2—applications

part 3—accelerationism

eprint.iacr.org/2020/107

One-shot Signatures and Applications to Hybrid
Quantum/Classical Authentication

Ryan Amos^{*1}, Marios Georgiou^{†2}, Aggelos Kiayias^{‡3}, and Mark Zhandry^{§4}

eprint.iacr.org/2020/107

revolutionary
and ignored

One-shot Signatures and Applications to Hybrid Quantum/Classical Authentication

Ryan Amos^{*1}, Marios Georgiou^{†2}, Aggelos Kiayias^{‡3}, and Mark Zhandry^{§4}

eprint.iacr.org/**2020/107**

revolutionary
and ignored

One-shot Signatures and Applications to Hybrid Quantum/Classical Authentication

Ryan Amos^{*1}, Marios Georgiou^{†2}, Aggelos Kiayias^{‡3}, and Mark Zhandry^{§4}



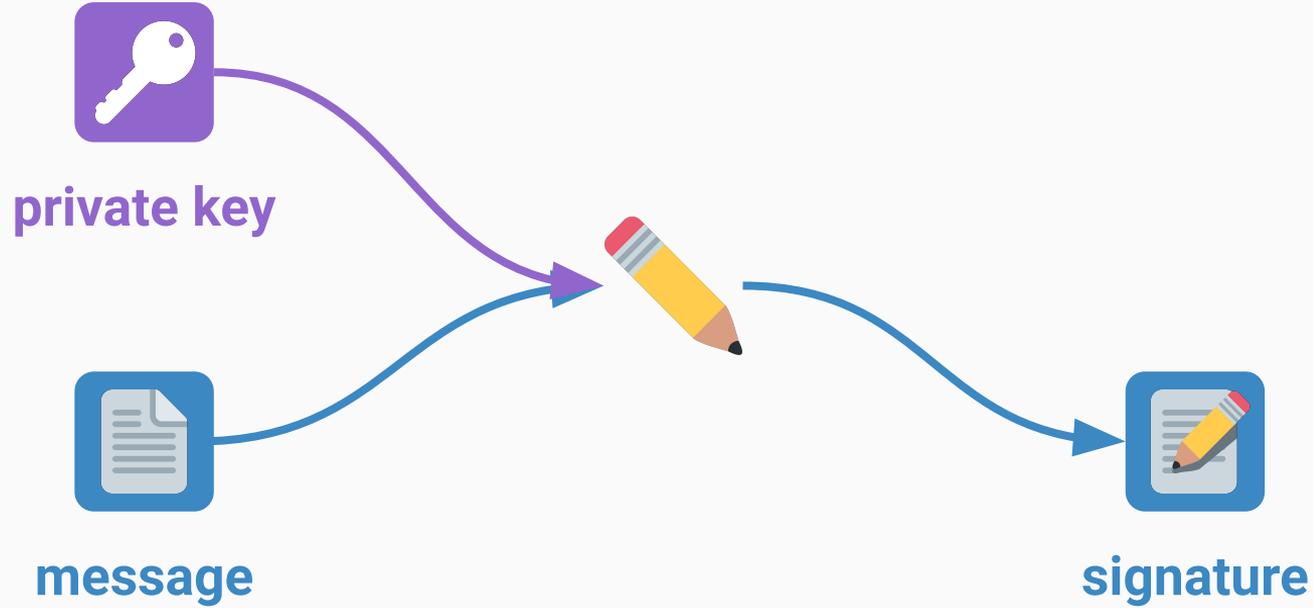


private key

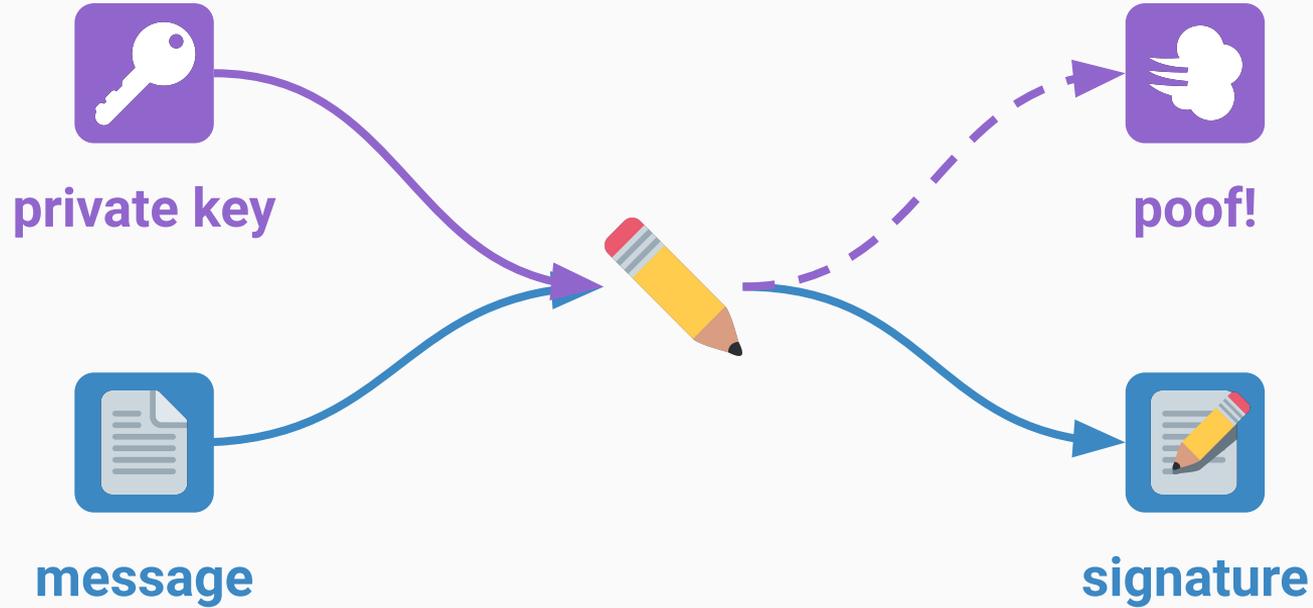


message

one-shot signing



one-shot signing



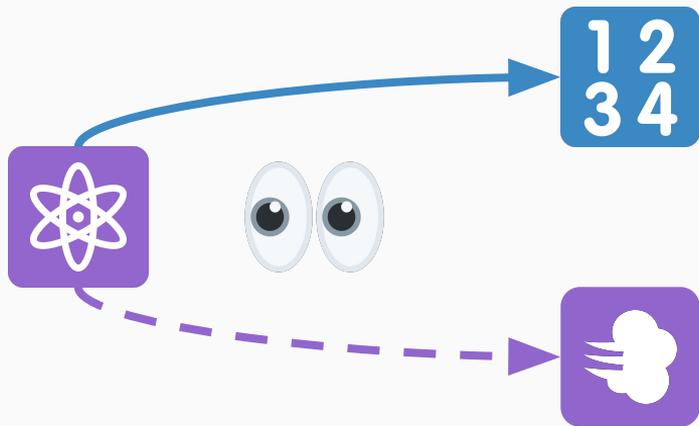
destructive measurements



no cloning



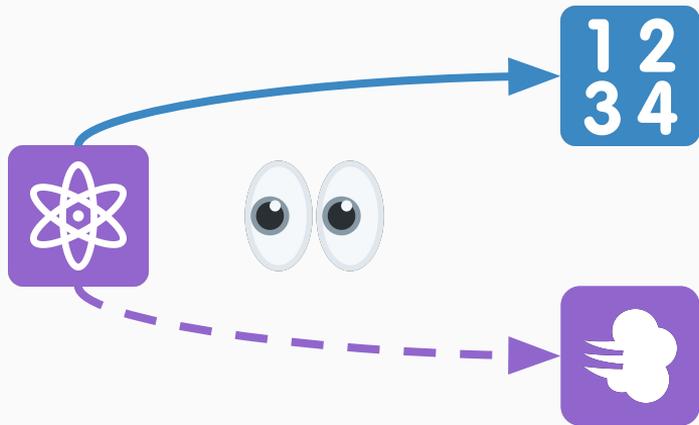
destructive measurements



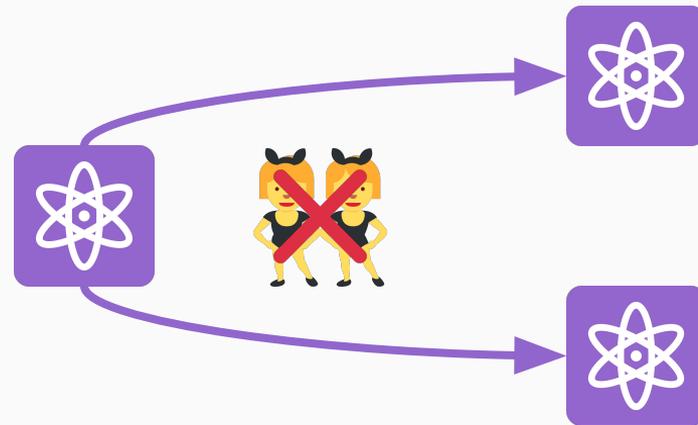
no cloning



destructive measurements



no cloning



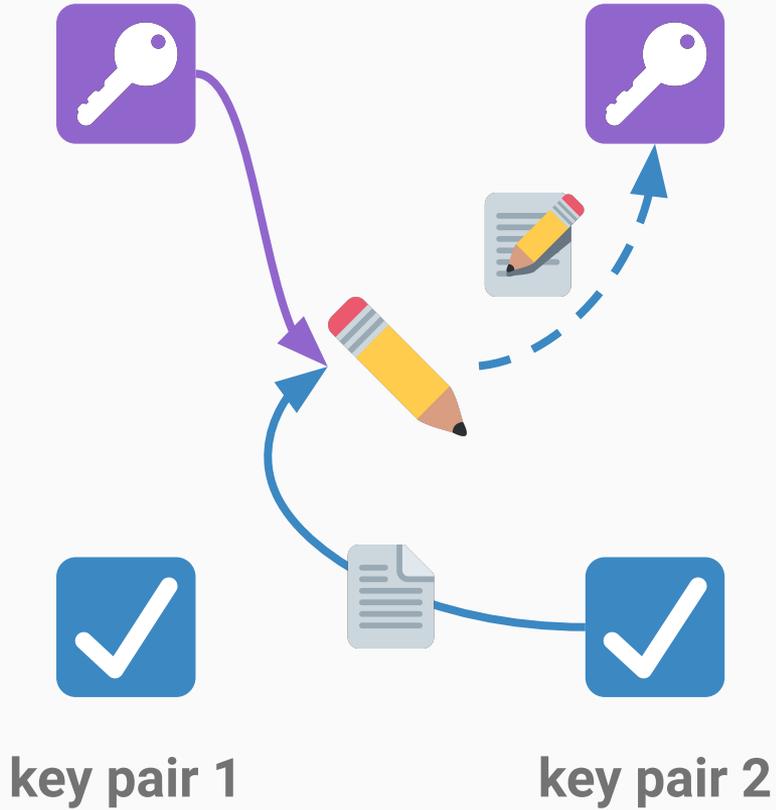
key chaining



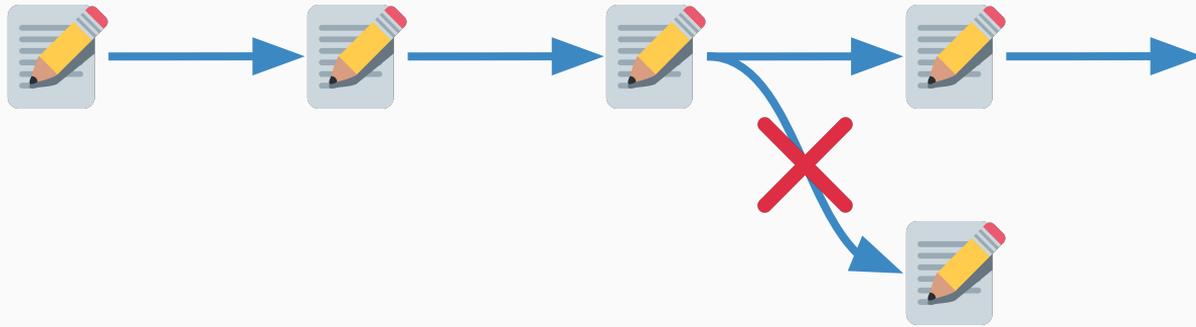
key pair 1

key pair 2

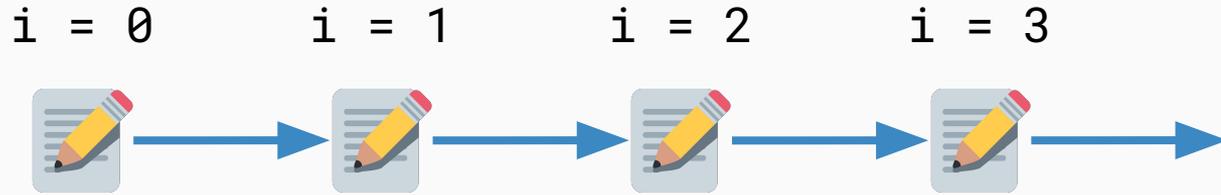
key chaining



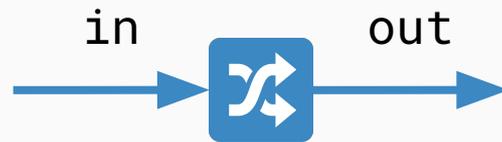
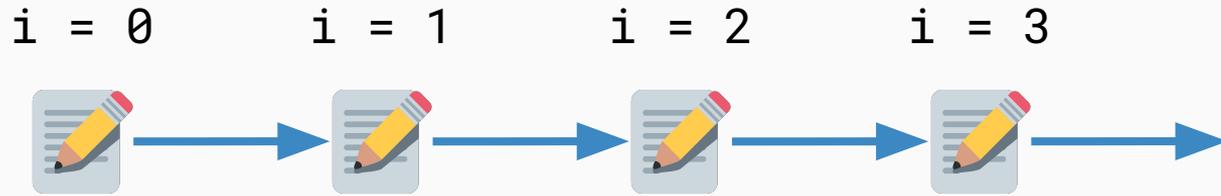
unforkable signature chains



stateful signature chains



stateful signature chains

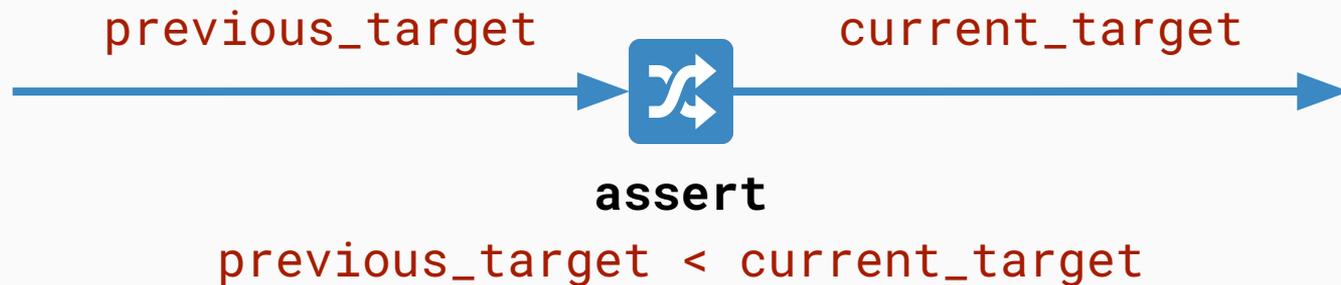


assert

`in + 1 = out`



no double vote



removing slashing conditions



no double vote



no surround vote

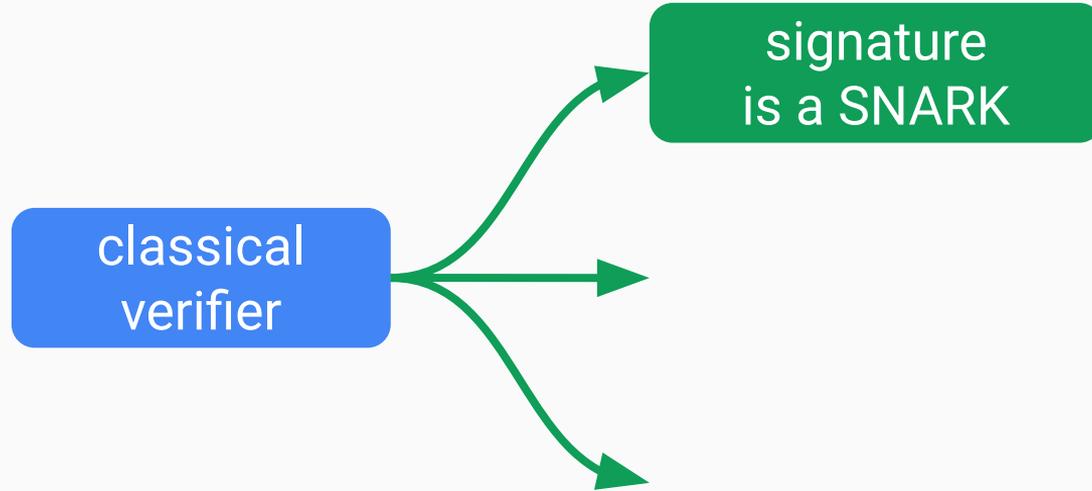
previous_source
previous_target

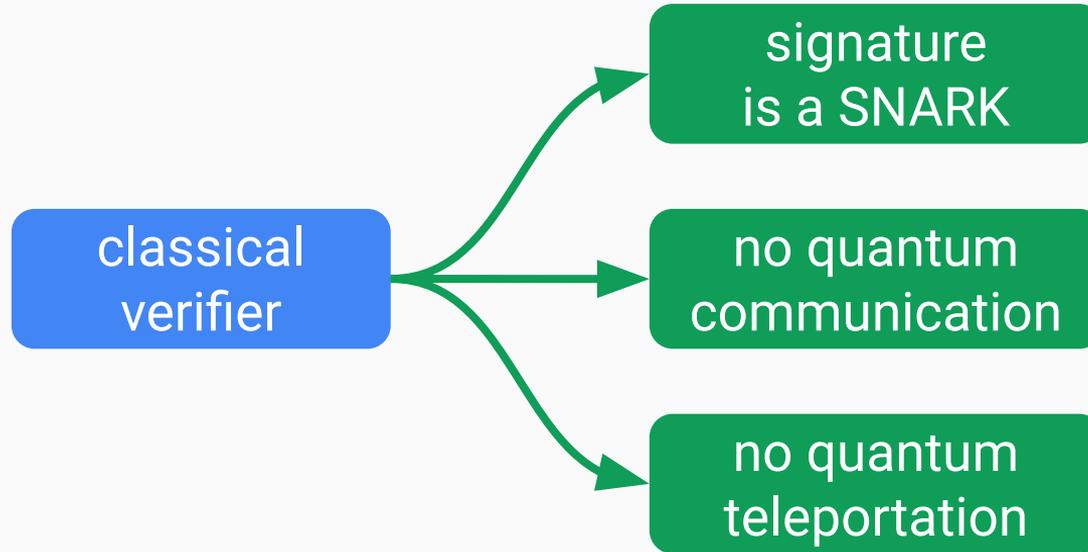
current_source
current_target



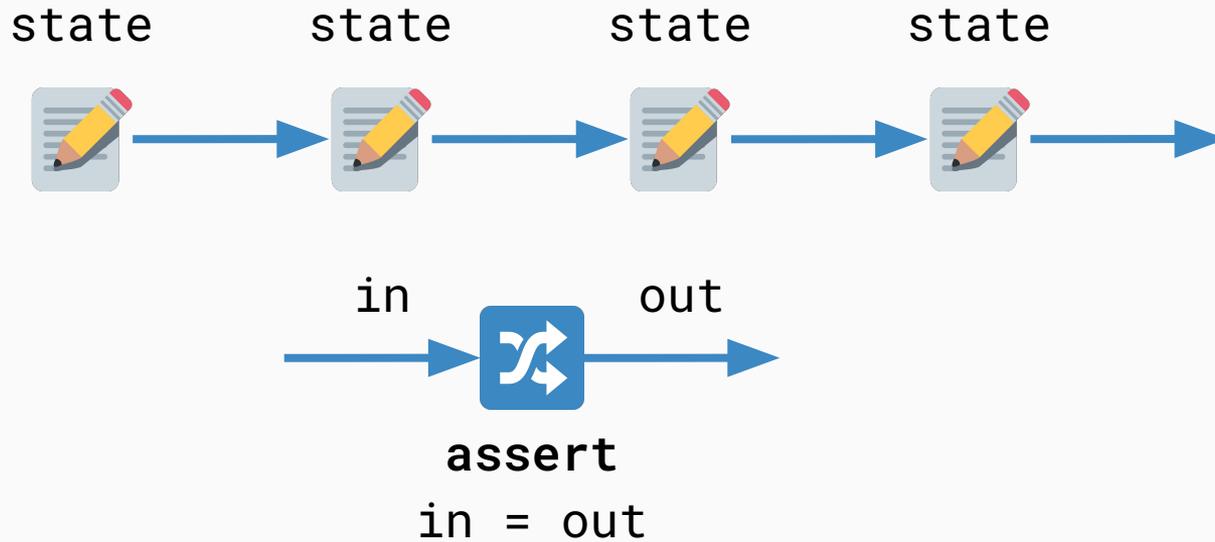
assert

previous_target < current_target
previous_source ≤ current_source

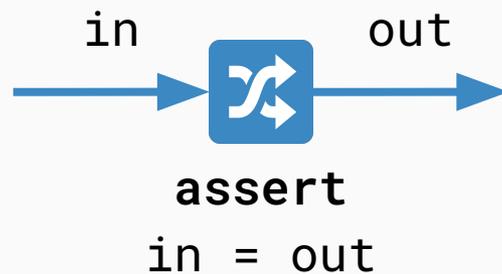
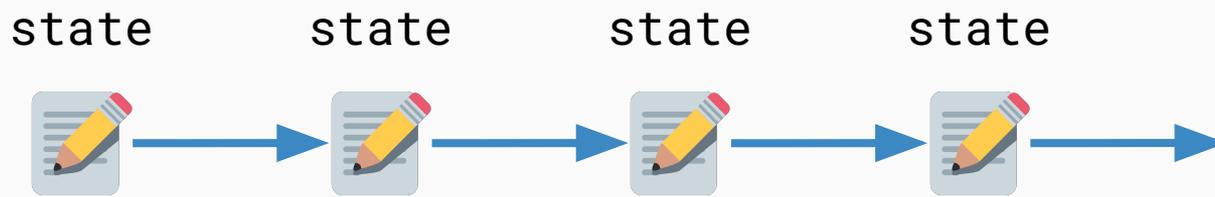




practicality—short-term memory



practicality—short-term memory

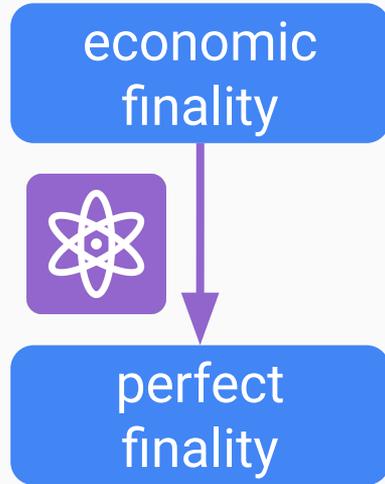


private key refresh

part 1—signature chains

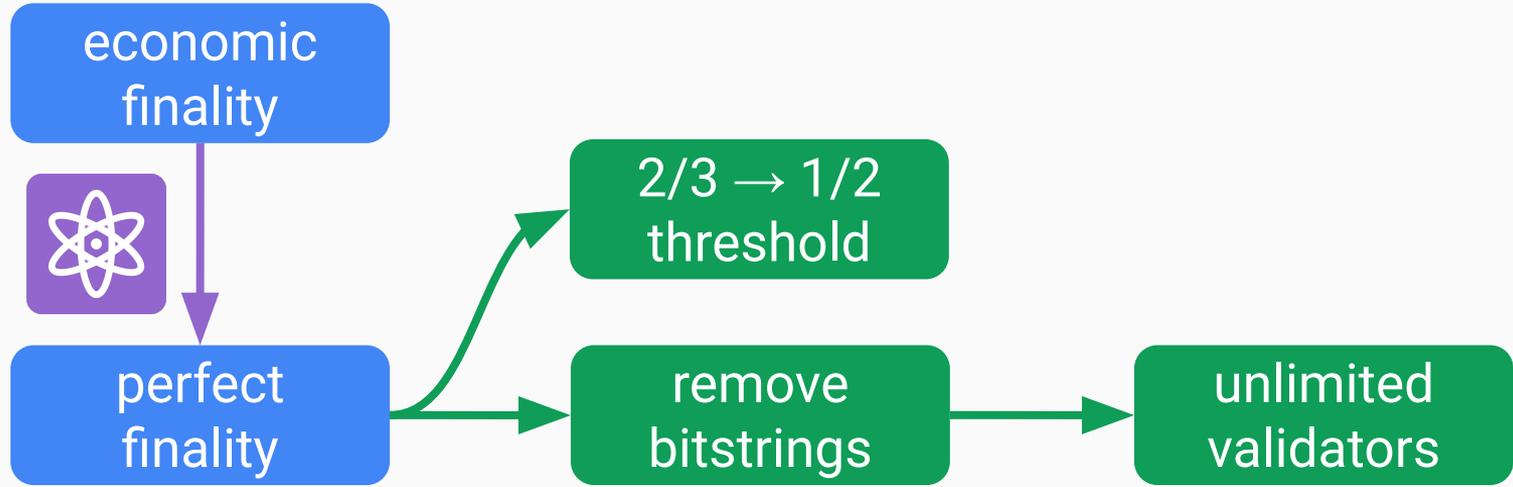
part 2—applications

part 3—accelerationism

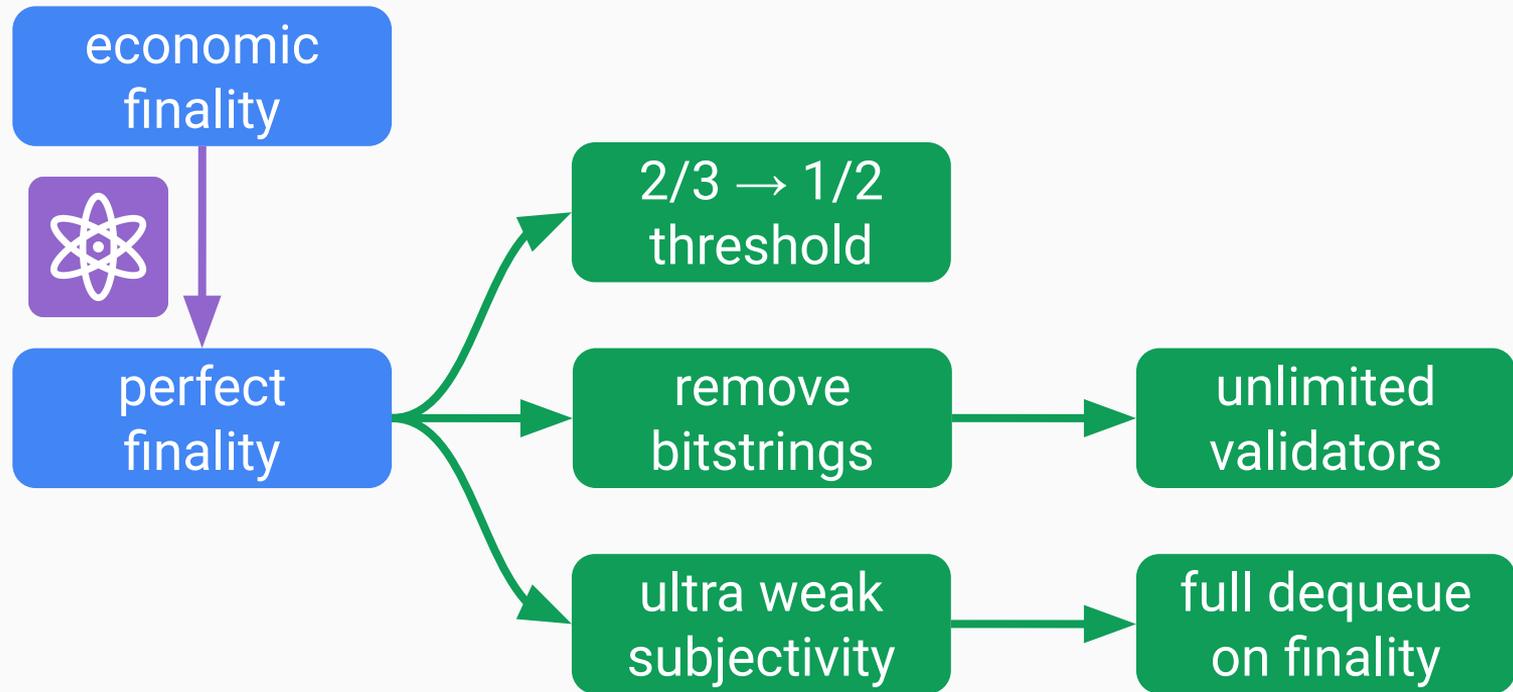




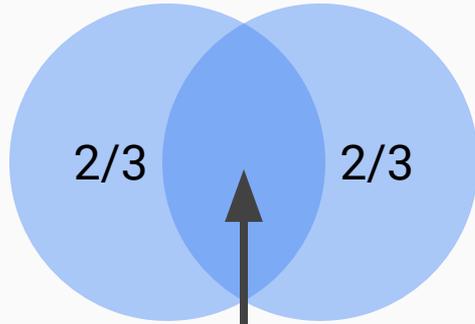
perfect finality



perfect finality



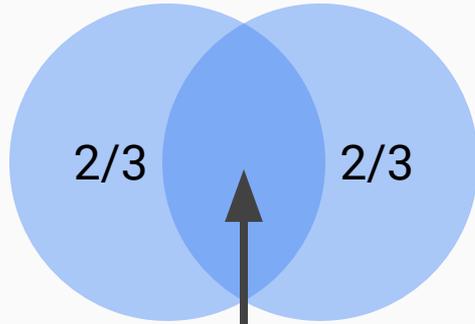
economic vs perfect finality



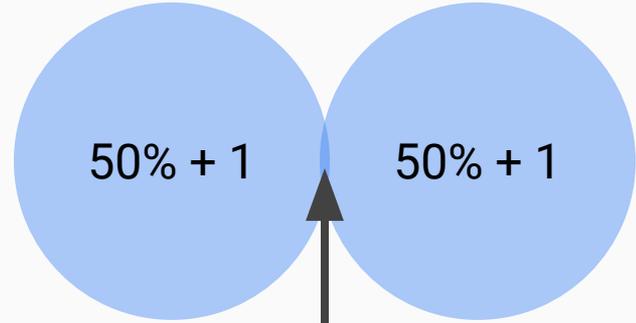
$\geq 1/3$
equivocations



economic vs perfect finality

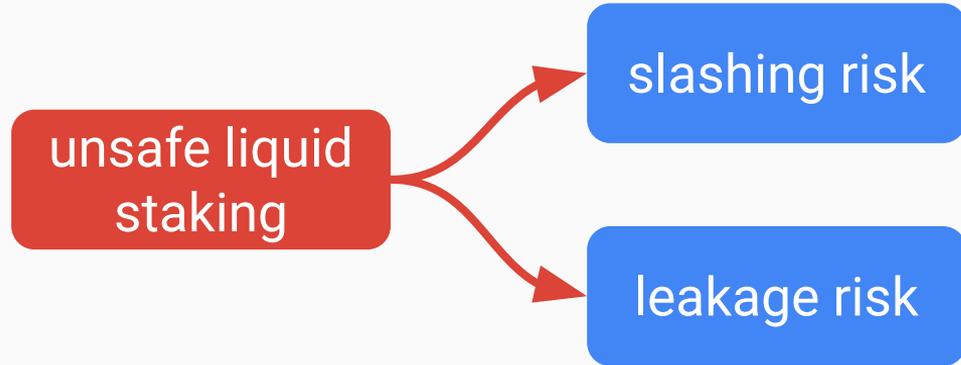


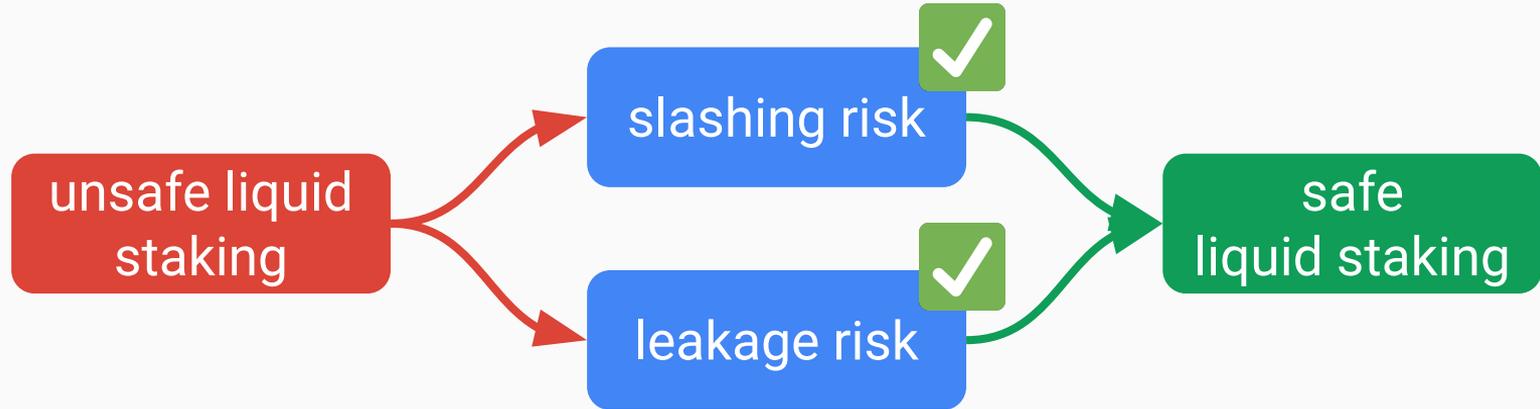
$\geq 1/3$
equivocations

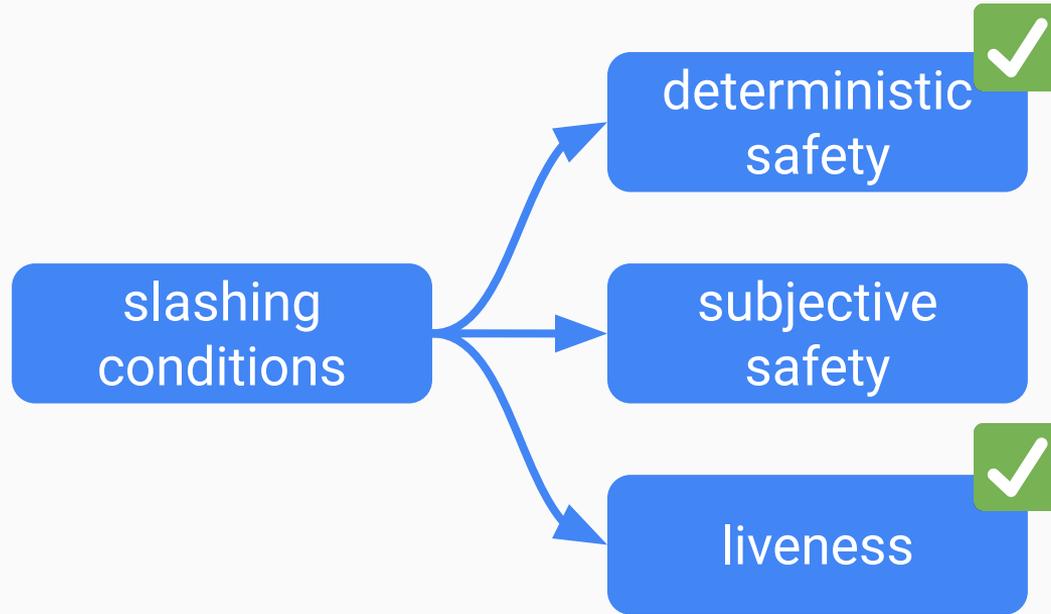


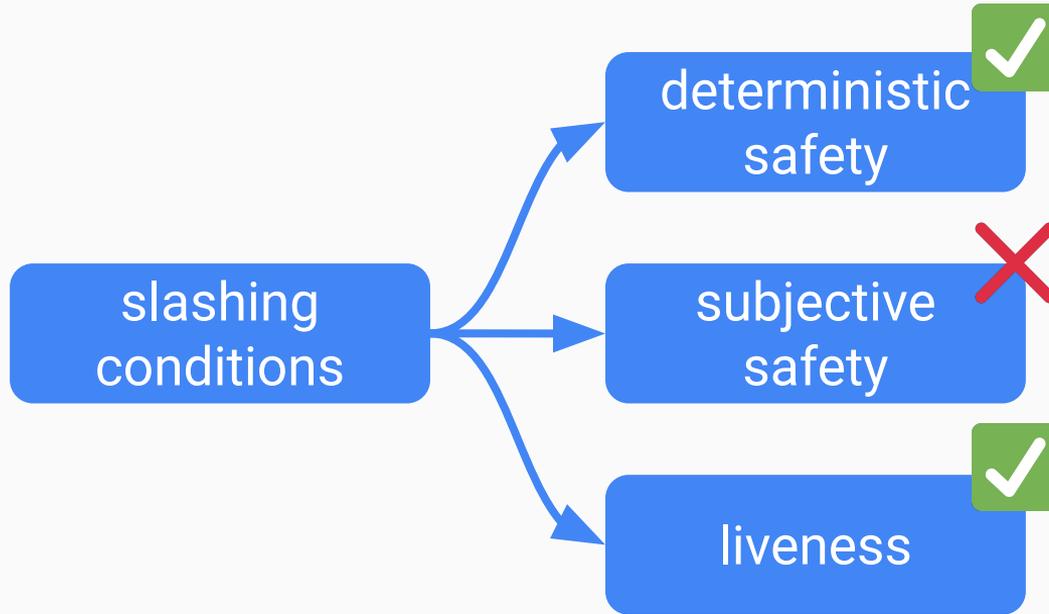
≥ 1
equivocations











quantum
money



quantum
money



proof of
location



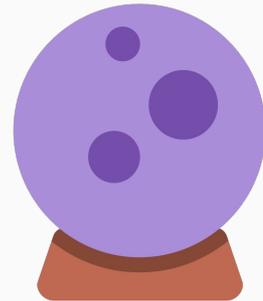
quantum
money



proof of
location



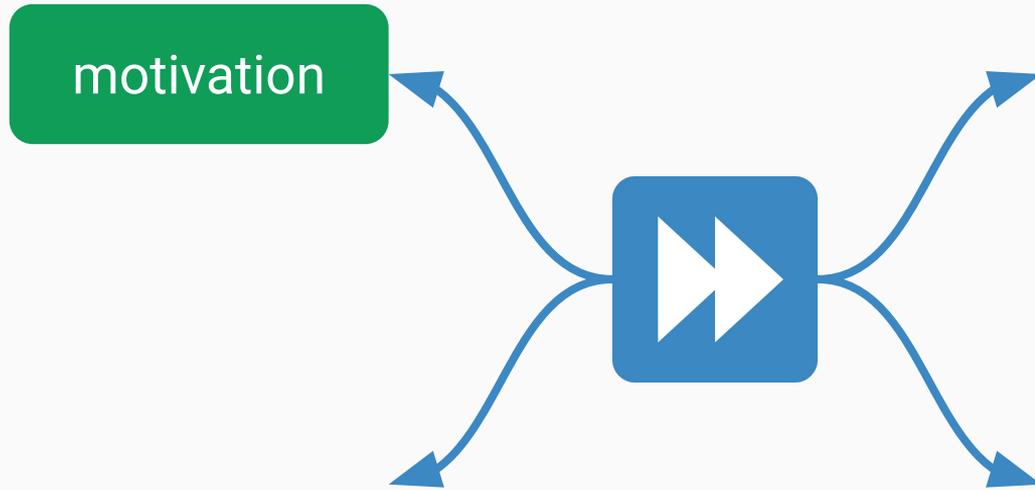
safer
oracles

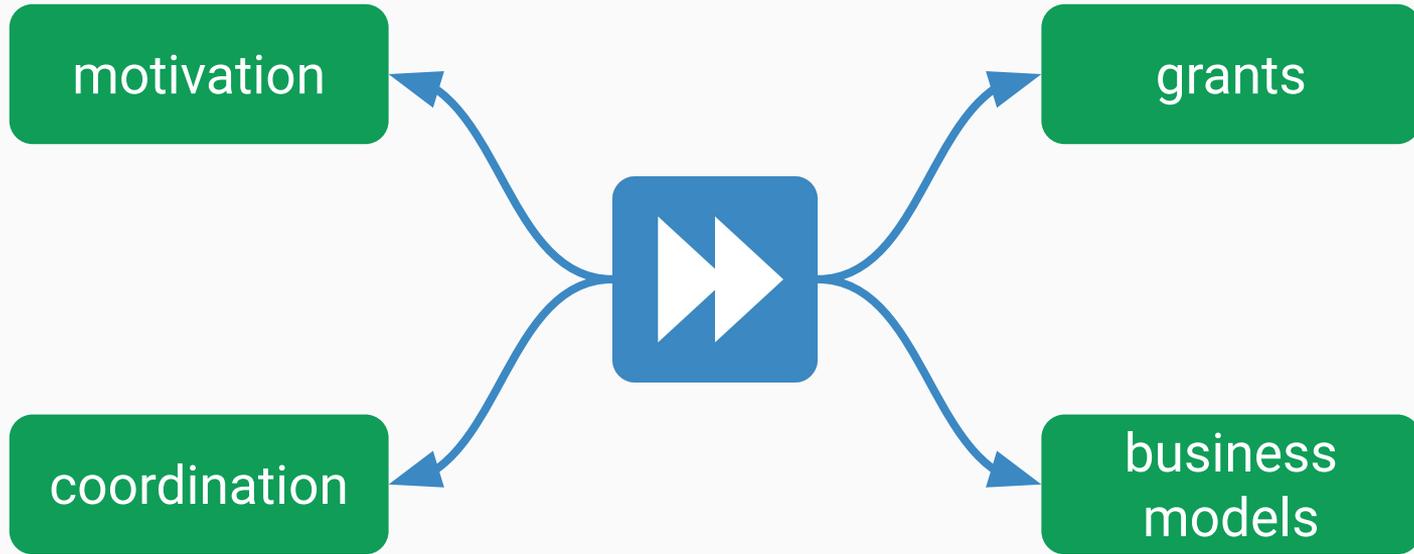


part 1—signature chains

part 2—applications

part 3—accelerationism







QSig



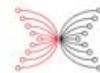
THE UNIVERSITY OF EDINBURGH
informatics



ethereum
foundation



Cardano
Foundation

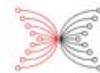


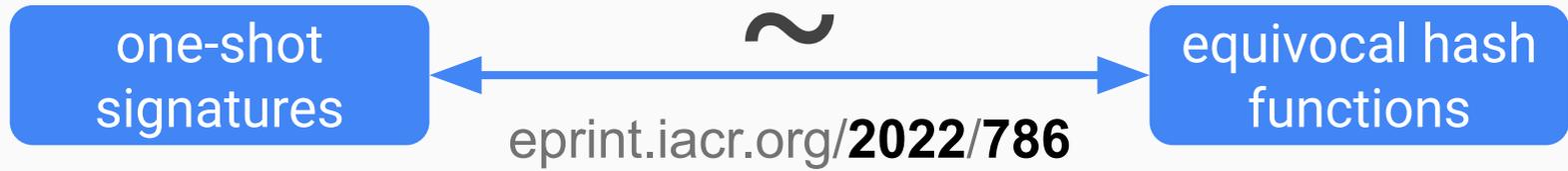
INPUT | OUTPUT

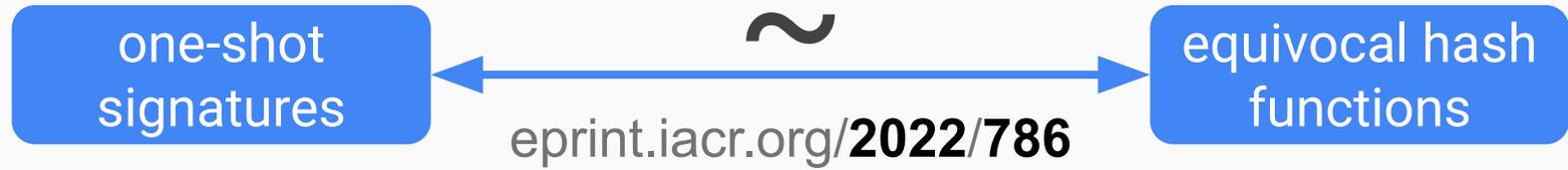


QSig

2025?



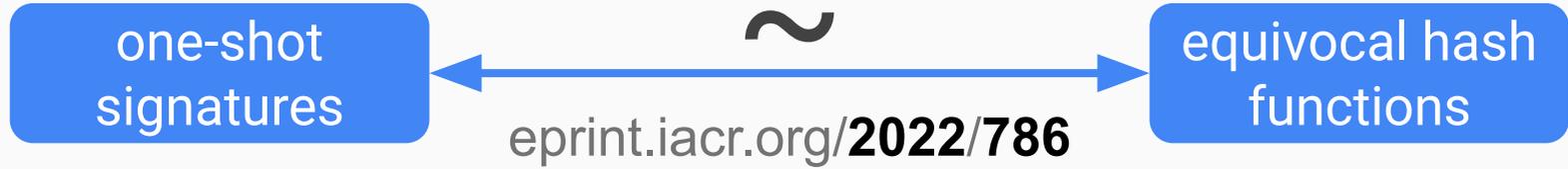




hash design
competition



hash function competition



hash design competition



NIST

SHA3, BLAKE



ethereum
foundation



STARKWARE

Poseidon, Rescue



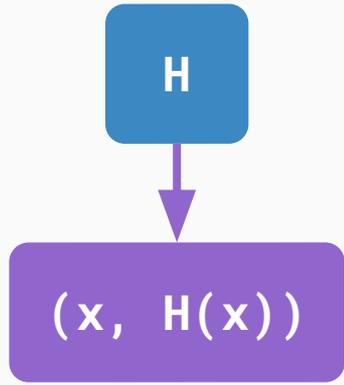
thank you :)

justin@ethereum.org

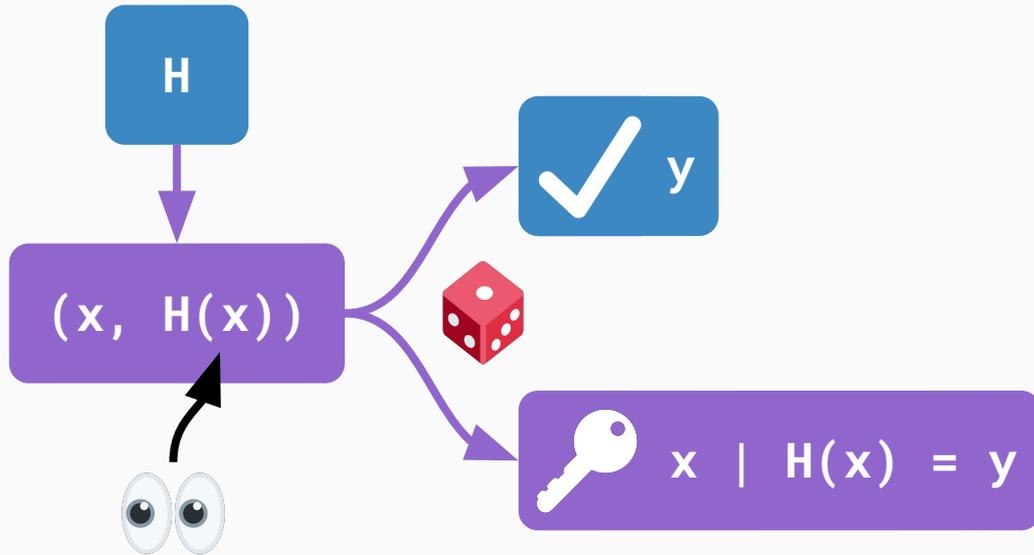
one quantum
computer

signature
reuse

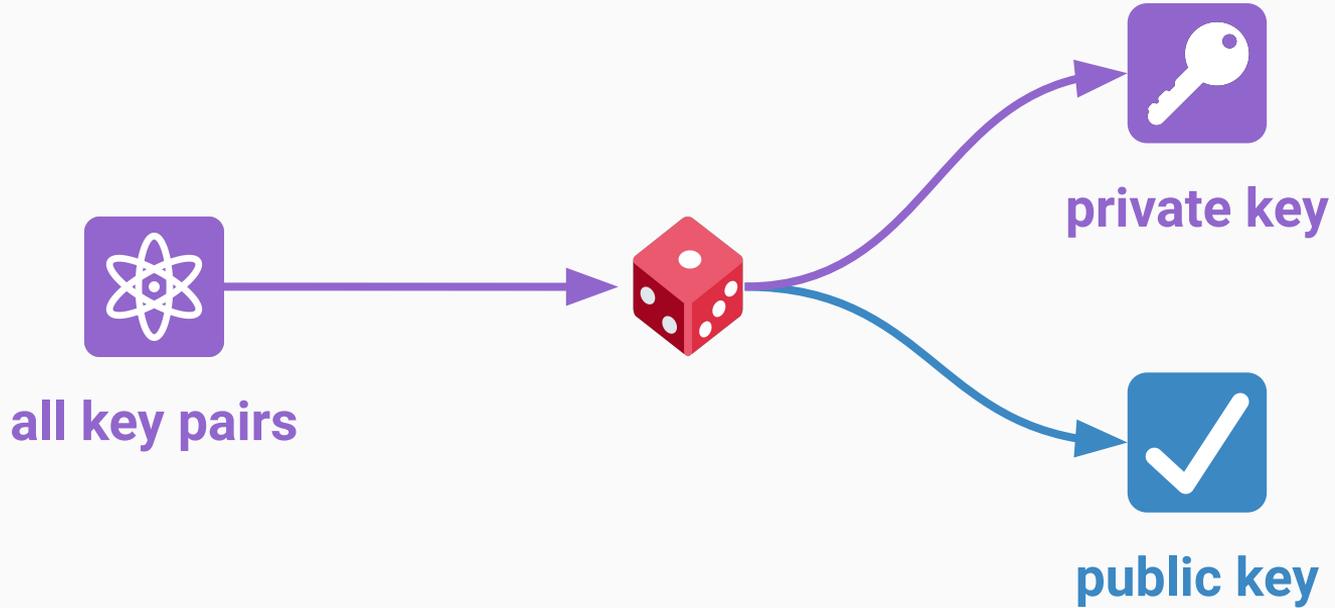
1-bit signing



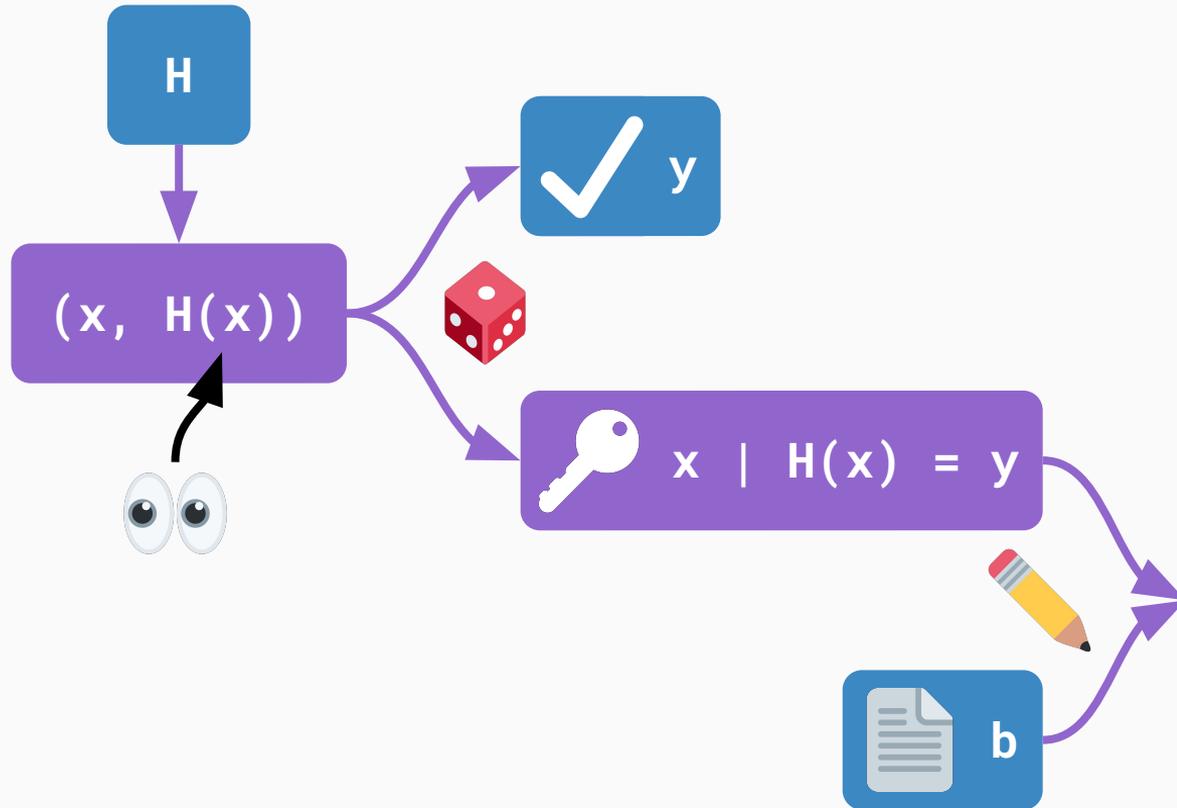
1-bit signing



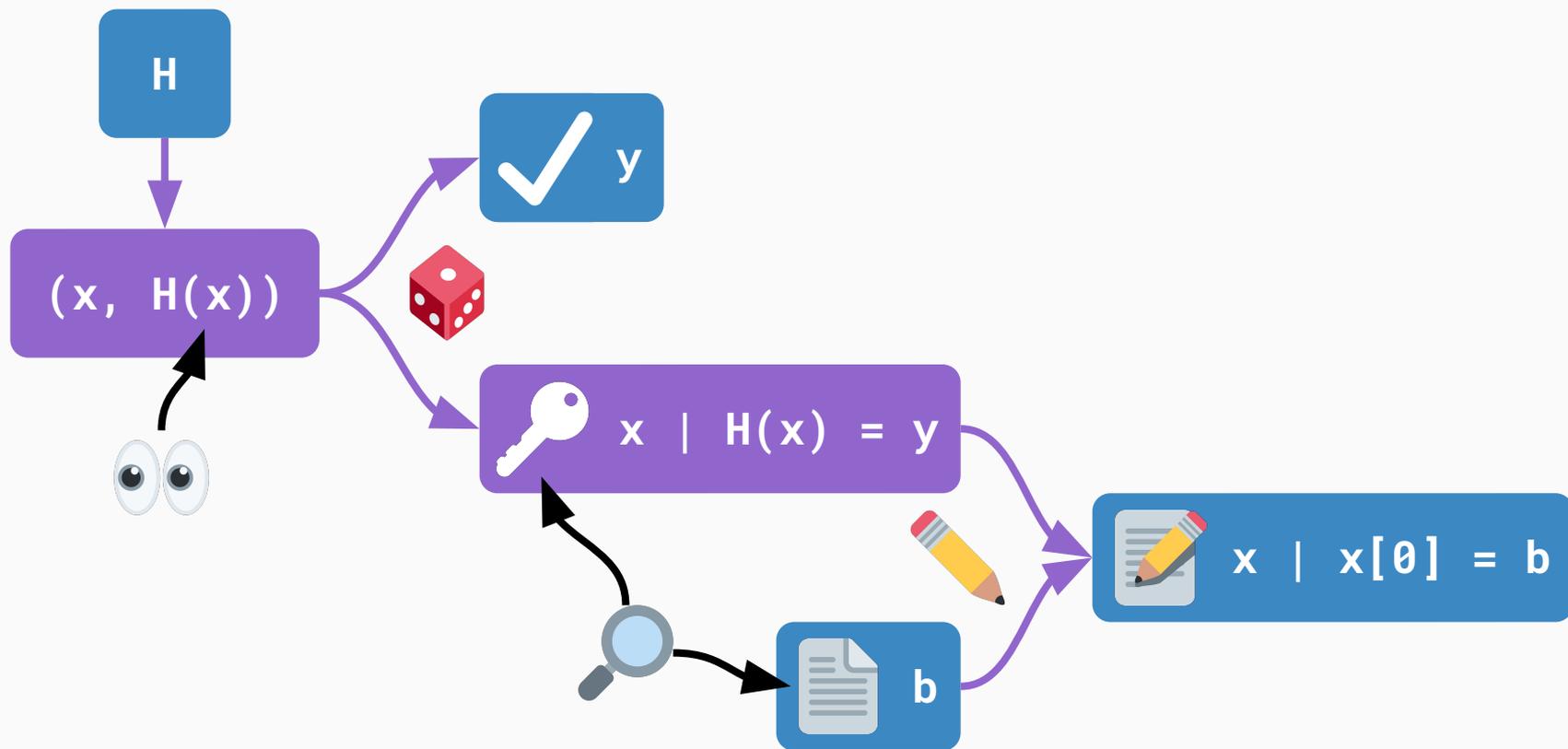
key generation



1-bit signing



1-bit signing



1. start with a suitable 256-bit hash function **H**

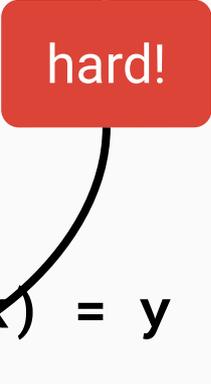
1. start with a suitable 256-bit hash function H
2. build a uniform superposition of all preimage-image pairs
 - $(\mathbf{x}, H(\mathbf{x}))$ for every 512-bit preimage \mathbf{x}

1. start with a suitable 256-bit hash function H
2. build a uniform superposition of all preimage-image pairs
 - $(\mathbf{x}, H(\mathbf{x}))$ for every 512-bit preimage \mathbf{x}
3. observe and collapse the second register to get
 - *pubkey*—a random image \mathbf{y}
 - *privkey*—a superposition of preimages \mathbf{x} such that $H(\mathbf{x}) = \mathbf{y}$

1. start with a suitable 256-bit hash function H
2. build a uniform superposition of all preimage-image pairs
 - $(\mathbf{x}, H(\mathbf{x}))$ for every 512-bit preimage \mathbf{x}
3. observe and collapse the second register to get
 - *pubkey*—a random image \mathbf{y}
 - *privkey*—a superposition of preimages \mathbf{x} such that $H(\mathbf{x}) = \mathbf{y}$
4. to sign bit \mathbf{b} run a fancy quantum search algorithm to find
 - *signature*—a preimage \mathbf{x} such that the first bit of \mathbf{x} is \mathbf{b}

1-bit signing

1. start with a suitable 256-bit hash function H
2. build a uniform superposition of all preimage-image pairs
 - $(\mathbf{x}, H(\mathbf{x}))$ for every 512-bit preimage \mathbf{x}
3. observe and collapse the second register to get
 - *pubkey*—a random image \mathbf{y}
 - *privkey*—a superposition of preimages \mathbf{x} such that $H(\mathbf{x}) = \mathbf{y}$
4. to sign bit \mathbf{b} run a fancy quantum search algorithm to find
 - *signature*—a preimage \mathbf{x} such that the first bit of \mathbf{x} is \mathbf{b}



hard!